





Advanced Computer Graphics Real-Time Rendering by Advanced Visibility Computations

G. Zachmann University of Bremen, Germany cgvr.informatik.uni-bremen.de



Remember the graphics pipeline



- A pipeline always has the throughput of its slowest link!
- Possible bottlenecks in the graphics pipeline :
 - In rasterizer → "fill limited"
 - In geometry stage → "transform limited"
 - Bus between app. and graphics hardware → "bus limited"
 - If the graphics card is faster than the application can provide geometry
 → "CPU limited" (recognizable by 100% CPU usage)

Classification of Visibility Problems

Bremen

W



- Problem classes within "visibility computations":
 - 1. Hidden Surface Elimination: which pixels (parts of polygons) are covered by others?
 - 2. Clipping: which pixels (parts of polygons) are inside the viewport?
 - **3.** Culling: which polygons cannot be visible? (e.g., because they are located behind the viewpoint)
- Difference: HSE & clipping are rather used to render an accurate image, culling is rather used to accelerate the rendering of large scenes
- Note: the boundary is blurred





- Let A = set of all primitives;
 let S = set of visible primitives.
- Many rendering algorithms operate on the entire set A, i.e., they have a minimum effort of O(|A|)
- No problem if $|S| \approx |A|$
- Also no problem, if the number of primitives is small compared to the number of pixels
 - Reminder: depth complexity

"to cull from" = "sammeln [aus ...] / auslesen" "to cull flowers" = Blumen pflücken





But for complex visual scenes, the number of visible primitives is typically much smaller than the total number of primitives!
 (i.e., |S| << |A|)





Culling is an important optimization technique (as opposed to clipping)





- For |S| << |A|, existing rendering algorithms are not efficient</p>
- Culling algorithms attempt to determine the set of non-visible primitives C = A \ S (or a subset thereof), or the set of visible primitives S (or superset thereof)
- Definition: potentially visible set (PVS) = a superset $S' \supseteq S$
 - Goal: compute PVS S' as small as possible, with minimal effort
 - Trivial PVS (with trivial effort) is, of course, A



Kinds of Culling







Back-Face Culling



- Definition: a solid = closed, opaque object = non-translucent object with non-degenerate volume
- Observations:
 - With solids, the back faces are never visible
 - For convex objects, there is exactly one *contiguous* back side
 - For non-convex solids, there may be several unconnected back sides





- Backface Culling = not drawing the surface parts that are on the far side, with respect to the viewpoint
 - Only works with solids!
- Compute normal n of the polygon
- Compute view vectors v from the viewpoint to all points p of the polygon
 - Perspevtive projection: v = p eye
 - Orthogonal projection: $\mathbf{v} = [0 \ 0 \ 1]^T$
- Polygon is back facing (don't draw), iff angle between n and v < 90°











$$N_{1} \cdot V = (2, 1, 2) \cdot (-1, 0, -1)$$

= -4 < 0

 $\Rightarrow N_{1}$ front facing
$$N_{2} \cdot V = (-3, 1, -2) \cdot (-1, 0, -1)$$

= 5 > 0

 $\Rightarrow N_{2}$ back facing
$$N_{2} = (-3, 1, -2)$$

 $N_{1} = (2, 1, 2)$
 $V = (-1, 0, -1)$



Backface Culling in OpenGL



Just enable it:

glCullFace(GL_BACK);
glEnable(GL_CULL_FACE);











- Central idea: replace the scalar product by classifying all normals
- Preprocessing: create classes over the set of all normals
 - Enclose the sphere of normals (aka. Gaussian sphere) with cube (direction cube)



- Results in $6 N^2$ classes (N = number of partitions along each axis)
- Classification of a normal is very easy
- With each polygon store the class of its normal





- Encoding a normal (pre-processing):
 - The entire direction cube \triangleq bit string of length $6 \cdot N^2$
 - A normal ≜ bit string with only one 1, otherwise 0
 - Encode this as offset + part of the bit string that contains the 1
 - E.g.: subdivide bit string in bytes, offset = 1 Byte, results in 256×8 = 2048 Bits

```
typedef struct PolygonNormalMask
{
    Byte offset, bitMask;
};
```

0....000001000000.....0 offset (in Bytes) bitMask

- Save those 2 bytes for each polygon
- E.g.: choose *N* = 16
- Results in 6.16.16 = 1536 classes for the set of all normals



Culling (initialization):

Bremen

llŰ

- Identify all those normal classes whose normals are all backfacing
- With orthographic projection:



 With perspective projection: which normals are backfacing depends on normal direction and position of the polygon!



 Therefore: determine a "conservative" set of classes which are backfacing – regardless of the location of the polygon







- In practice:
 - Test each class in all four corners of the view frustum
 - Test for a class = test of 4 normals, which are pointing to the corners of the cell (on the direction cube) that represents that class





Represent this conservative set of classes as a bit string (e.g. 2048
 Bits = 256 Bytes) in a byte array:

Byte BackMask[256];

Culling (runtime): test for each polygon

```
if ( (BackMask[byteOffset] & polygon.bitMask) == 1 )
    render polygon
```

- Further acceleration:
 - Divide view frustum into sectors
 - Render the scene "sector by sector"
 - Thus, the angle $\alpha/2$ in each sector is smaller
 - For each sector, compute its own BackMask[]





216 classes ("clusters")

1536 classes ("clusters")



BackMask for the current viewpoint (green = backfacing)









Result: speedup factor ~1.5 compared to OpenGL backface culling





Reminder: some simple rules for min/max

$$\max_{i} \{x_{i} + y_{i}\} \leq \max_{i} \{x_{i}\} + \max_{i} \{y_{i}\}$$
$$\max_{i} \{x_{i} - y_{i}\} \leq \max_{i} \{x_{i}\} - \min_{i} \{y_{i}\}$$
$$\max_{i} \{kx_{i}\} = \begin{cases} k \max_{i} \{x_{i}\} & , \ k \geq 0 \\ k \min_{i} \{x_{i}\} & , \ k < 0 \end{cases}$$

- In the following, nⁱ and pⁱ are the normal and a vertex of a polygon from a cluster (a set) of polygons; let e be the viewpoint
- Attention: in the following, we use the "inverted" definition for backfacing!

$$\mathbf{n} \cdot (\mathbf{e} - \mathbf{p}) \le 0$$





- Assumption: cluster (= set) of polygons is given
- All polygons in cluster are backfacing if and only if

$$\forall i : \mathbf{n}^{i} \left(\mathbf{e} - \mathbf{p}^{i} \right) \leq 0 \quad \Leftrightarrow \\ \max \left\{ \mathbf{n}^{i} \left(\mathbf{e} - \mathbf{p}^{i} \right) \right\} \leq 0 \tag{1}$$

Upper bound for (1) is

$$\max\left\{\mathbf{n}^{i}\left(\mathbf{e}-\mathbf{p}^{i}\right)\right\} \leq \max\left\{\mathbf{en}^{i}\right\} - \min\left\{\mathbf{n}^{i}\mathbf{p}^{i}\right\}$$
(2)

- Set $d := \min\{n^{i}p^{i}\}$ (pre-computation)
- Write (2) as

$$\max \left\{ \mathbf{n}^{i} \left(\mathbf{e} - \mathbf{p}^{i} \right) \right\} \leq \max \left\{ e_{x} n_{x}^{i} + e_{y} n_{y}^{i} + e_{z} n_{z}^{i} \right\} - d$$
$$\leq \max \left\{ e_{x} n_{x}^{i} \right\} + \max \left\{ e_{y} n_{y}^{i} \right\} + \max \left\{ e_{z} n_{z}^{i} \right\} - d \quad (3)$$





Assumption: e is located in the positive octant, i.e., e_x, e_y, e_z ≥ 0; then we can give an upper bound of (3):

 $\max \left\{ \mathbf{n}^{i} \left(\mathbf{e} - \mathbf{p}^{i} \right) \right\}$ $\leq e_{x} \cdot \max\{n_{x}^{i}\} + e_{y} \cdot \max\{n_{y}^{i}\} + e_{z} \cdot \max\{n_{z}^{i}\} - d$ $\leq \mathbf{m} \cdot \mathbf{e} - d , \quad \text{mit} \quad \mathbf{m} = \begin{pmatrix} \max\{n_{x}^{i}\} \\ \max\{n_{y}^{i}\} \\ \max\{n_{z}^{i}\} \end{pmatrix}$

• Analogously, for e_x , e_y , $e_z \le 0$:

$$\max\left\{\mathbf{n}^{i}\left(\mathbf{e}-\mathbf{p}^{i}\right)\right\} \leq \mathbf{m}' \cdot \mathbf{e} - d , \quad \text{with} \quad \mathbf{m}' = \begin{pmatrix}\min\{n_{x}^{i}\}\\\min\{n_{y}^{i}\}\\\min\{n_{z}^{i}\}\end{pmatrix}$$





- For all other octants, combine min and max appropriately
- We can write this with kind of a "combination" operator on vectors

$$\mathsf{comb}(\mathbf{u},\mathbf{v};\mathbf{e}) := \mathbf{w} \quad \mathsf{with} \quad w_{\alpha} = \begin{cases} u_{\alpha} & , e_{\alpha} \leq 0 \\ v_{\alpha} & , e_{\alpha} > 0 \end{cases}$$
, $\alpha \in \{x, y, z\}$

• This allows us write the (conservative) test as:

$$\operatorname{comb}(\mathbf{m}', \mathbf{m}; \mathbf{e}) \cdot \mathbf{e} - d \leq 0 \implies \operatorname{cluster} \operatorname{is} \operatorname{backfacing}$$
 (4)

- Pre-computation: for each cluster determine m, m' and d
- Memory requirements per cluster: 28 Bytes (2 vectors + 1 scalar)



Bremen

UŰ



- Inequality (4) defines 8 planes (one per octant)
- The 4 planes of adjacent octants intersect at one point, which lies on the coordinate axis "between" the 4 octants
 - Example: Consider the 4 planes in the octants with $e_X \ge 0$
 - All 4 planes have normals of the form $n = (m_x, \cdot, \cdot)$
 - So, they all intersect the x-axis at the point $\left(\frac{d}{m_x}, 0, 0\right)$.
- Those 8 planes form a closed volume, the so-called culling volume







- Problem: if the polygons are far away from the origin, and the origin is located on the positive side of the normal, then d is very much negative \rightarrow the test is never positive
- Solution: run the test in a local coordinate system by
- Move the local origin c such that

$$d = \min\left\{\mathbf{n}^i \cdot \left(\mathbf{p}^i - \mathbf{c}\right)\right\}$$

is as large (and positive) as possible

- Wanted is the optimal c
 - Question: Will rotation achieve something?
 - In practice: Try the center and corner of the BBox of the cluster as c
- Save c with the cluster, then test

$$\operatorname{comb}(\mathbf{m}',\mathbf{m};\mathbf{e}-\mathbf{c})\cdot(\mathbf{e}-\mathbf{c})-d\leq 0$$





Two clusters can be combined to form a joint cluster:

Bremen

$$\hat{\mathbf{m}'} = \begin{pmatrix} \min(m'_x^1, m'_x^2) \\ \min(m'_y^1, m'_y^2) \\ \min(m'_z^1, m'_z^2) \end{pmatrix} \quad \hat{\mathbf{m}} = \begin{pmatrix} \max(m_x^1, m_x^2) \\ \max(m_y^1, m_y^2) \\ \max(m_z^1, m_z^2) \end{pmatrix} \\ \hat{d} = \min(d_1, d_2)$$

- These two vectors and \hat{d} provide a conservative estimate
- I.e.: if the joint cluster is invisible, then the two original clusters are guaranteed to be invisible, too → cluster hierarchy
- If a hierarchy of clusters is created, define a front-facing test, analogously to the backfacing test:
 - Stop testing, if a complete joint cluster is front- or back-facing
 - Otherwise: test the children for being completely front- or back-facing



Generating the Clusters



- For the evaluation of cluster candidates in an algorithm, we need a measure of the "performance" of a cluster
- Here: probability P that the cluster will be culled
- Use a heuristic to calculate *P* :

$$\frac{\text{Vol(culling volume)}}{\text{Vol(all possible viewpoint position)}} = \frac{\text{Vol}(C)}{\text{Vol}(U)}$$

- Vol(C) can be computed exactly
- For U choose the BBox of the entire scene
- If local culling coordinates are used: choose U = c · Bbox(cluster) ("near-culling probability")







• Question: given two clusters A , B; Is it faster to test and to render A and B separately, or is it faster to test the joint cluster $C = A \cup B$ first? (on average!)

Let T(A) be the expected(!) time to test cluster A and render it in case of (possible) visibility. Then

T(A) = t + (1 - P(A)) R(A)

where P(A) = probability, that cluster A gets culled, R(A) = time to render A (without further tests), and t = time for backface test of a cluster





• So, combining clusters A and B is worth it, if and only if

$$T(C) < T(A) + T(B)$$
 \Leftrightarrow

$$t+\left(1-P(C)
ight)R(C)<2t+\left(1-P(A)
ight)R(A)+\left(1-P(B)
ight)R(B)$$

$$P(C) > rac{-t + P(A)R(A) + P(B)R(B)}{R(A) + R(B)} \quad \Leftrightarrow$$

$$P(C) > \frac{P(A)n_A + P(B)n_B - \frac{t}{r}}{n_A + n_B}$$
Assumption:

$$R(A) = n_A \cdot r,$$

$$r = \text{constant effort}$$
for one polygon

 Ratio t/r depends on the machine; but can easily be determined experimentally and automatically in advance (depends on graphics card, number of light sources, textures, ...)